

Continuity

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1 Real numbers

The real numbers consist of: a set \mathbf{R} , together with two binary operations, formally functions $\mathbf{R} \times \mathbf{R} \rightarrow \mathbf{R}$, written $(a, b) \mapsto a + b$ and $(a, b) \mapsto ab$; two distinguished elements (i.e. elements of \mathbf{R} with definite names) 0 and 1; and a binary relation written $a < b$. This set-up obeys the usual axioms (associativity etc.).

2 The real numbers as a complete ordered field

Definition 2.1 *A system satisfying the usual axioms of addition and multiplication for numbers is called an (algebraic) field.*

Definition 2.2 *A system satisfying the usual axioms of addition and multiplication for numbers, and the usual axioms for ordering, is called an ordered field.*

Completeness axiom: *Every non-empty bounded subset S of \mathbf{R} has a least upper bound $\sup S$ in \mathbf{R} .*

Suppose now that \mathbf{F} is (just known to be) an ordered field. Then the Completeness axiom makes sense in \mathbf{F} , by the following definitions:

Definition 2.3 *For $x, y \in \mathbf{F}$ we define $x \leq y$ to mean (either $x < y$ or $x = y$).*

Definition 2.4 *If $S \subseteq \mathbf{F}$ is a subset then $U \in \mathbf{F}$ is called an upper bound for S if $s \leq U \forall s \in S$.*

Definition 2.5 *If $S \subseteq \mathbf{F}$ ($S \neq \emptyset$) has an upper bound then we say that S is bounded above.*

Then the Completeness axiom makes sense in \mathbf{F} , though it may or may not be true in \mathbf{F} . We assume it is true for \mathbf{R} , so we take \mathbf{R} to be a complete ordered field.

Definition 2.6 *A least upper bound ℓ for a nonempty subset S of any ordered field \mathbf{F} is an element of \mathbf{F} such that:*

1. ℓ is an upper bound for S , and
2. If u is any upper bound for S then $\ell \leq u$.

Traditionally we write $\sup S$ for the least upper bound of $S \subseteq \mathbf{R}$.

We'll assume \mathbf{R} contains subsets $\mathbf{N}, \mathbf{Z}, \mathbf{Q}$ with all the usual properties.

Proposition 2.7 *If $S \subseteq \mathbf{R}$ has a least upper bound, then it has at most one.*

Proof: Suppose ℓ and ℓ' are both least upper bounds for $S \subseteq \mathbf{R}$. View ℓ as *an* upper bound and ℓ' as *least* such, and see from (2) that $\ell' \leq \ell$. Now interchange the rôles of ℓ and ℓ' , and get $\ell \leq \ell'$. So $\ell = \ell'$.

Since by the axiom any non-empty $S \subseteq \mathbf{R}$ which is bounded above has a least upper bound, we now see that it is unique. Call it $\sup S$ (supremum or l.u.b.).

2.1 Remark

If $F \subseteq \mathbf{R}$ is finite (non-empty) then F contains a *greatest* member, and this is $\sup F$.

2.2 Remark

More generally (even if S isn't finite) it may be that S contains a greatest element; call it x_{\max} . Check:

1. For any $x \in S$, $x \leq x_{\max}$, by definition of greatest element.
2. if U is any bound for S , then (since $x_{\max} \in S$), we get $x_{\max} \leq U$.

Examples

1. In $[0, 1]$, $x_{\max} = 1$.
2. In $(0, 1)$, there's no greatest element.
3. In $\{1 - \frac{1}{n} : n \in \mathbf{N}_+\} = \{0, \frac{1}{2}, \frac{2}{3}, \dots\}$, there's no greatest element.

NB: Not all subsets of \mathbf{R} are intervals.

2.3 Remark

If $S \subseteq \mathbf{R}$ is non-empty and bounded above but has no greatest element, then $\sup S$ is a kind of substitute for the greatest element x_{\max} .

Proposition 2.8 (The Archimedian Property) .

- (a) For any $x \in \mathbf{R}$, $\exists n \in \mathbf{N}$ such that $n > x$.
- (b) For any $x > 0$ in \mathbf{R} , $\exists n \in \mathbf{N}$ such that $\frac{1}{n} < x$.

Proof

- (a) Suppose there is some $x \in \mathbf{R}$ for which (a) fails, so $n \leq x \forall n \in \mathbf{N}$.
So \mathbf{N} is non-empty and bounded above. By the completeness axiom, \mathbf{N} has a supremum, say ℓ .
For any $n \in \mathbf{N}$, we have $n + 1 \in \mathbf{N}$, so $n + 1 \leq \ell$. Hence $n < \ell - 1$.
But this is true for any $n \in \mathbf{N}$, so $\ell - 1$ is another upper bound for \mathbf{N} , so ℓ is not the least such. This contradicts the choice of ℓ , so (a) must be true.
- (b) If $x > 0$ then $\frac{1}{x} > 0$, so by (a) $\exists n \in \mathbf{N}$ with $n > \frac{1}{x}$, so $\frac{1}{n} < x$.

2.4 Remark

For some subsets of \mathbf{R} , even when there's no greatest element, we can see that a supremum exists without having to use the completeness axiom.

Example

To prove that $\sup(0, 1) = 1$: let $S = (0, 1)$.

1. $x \in S \Rightarrow x < 1$, by definition of $(0, 1)$.
2. Let U be any upper bound for S . We need to show that $1 \leq U$. Suppose for a contradiction that $U < 1$. So $1 - U > 0$. So by 1.8(b), $\exists n \in \mathbf{N}$ such that $1 - U > \frac{1}{n}$. So $1 - \frac{1}{n} > U$. But $1 - \frac{1}{n} \in (0, 1)$ (NB: we can w.l.o.g.¹ assume that $n > 1$). So U is not an upper bound for S .

2.5 Remark

We now look at an example where we *do* need the axiom.

Consider $S = \{x \in \mathbf{R} : x^2 < 2\}$. Then $S \neq \emptyset$ (e.g. $1 \in S$) and S is bounded above, e.g. by 2, since if $x > 2$ then $x^2 > 4$, so $x \notin S$.

By completeness, S has a supremum, say ℓ . We'll show that $\ell^2 = 2$, i.e. $\ell = \sqrt{2}$.

Proof (by contradiction)

1. Suppose $\ell^2 < 2$. We'll show that ℓ is not an upper bound for S , by showing that $\exists n \in \mathbf{N}$ with $\ell + \frac{1}{n} \in S$.

We want $\left(\ell + \frac{1}{n}\right)^2 < 2$, i.e. $\ell^2 + \frac{2\ell}{n} + \frac{1}{n^2} < 2$.

Equivalently $\frac{2\ell}{n} + \frac{1}{n^2} < 2 - \ell^2$ (NB: $2 - \ell^2 > 0$ by assumption).

Well, $\frac{2\ell}{n} + \frac{1}{n^2} \leq \frac{2\ell}{n} + \frac{n}{n^2}$, so enough to get n such that $\frac{2\ell + 1}{n} < 2 - \ell^2$.

But by 1.8(a), $\exists n \in \mathbf{N}$ with $n > \frac{2\ell + 1}{2 - \ell^2}$.

2. Suppose $\ell^2 > 2$. We'll show that ℓ isn't "least" by finding n such that $\ell - \frac{1}{n}$ is an upper bound of S .

Enough to see $\left(\ell - \frac{1}{n}\right)^2 > 2$ for then if $x \in S$, we have:

$x^2 < 2 < \left(\ell - \frac{1}{n}\right)^2$ and we get $x < \ell - \frac{1}{n}$.

(NB: $\ell - \frac{1}{n} > 0$ since $1 \in S$, so $\ell \geq 1$)

We want $\ell^2 - \frac{2\ell}{n} + \frac{1}{n^2} > 2$. Enough to get $\ell^2 - \frac{2\ell}{n} > 2$

¹without loss of generality

, or equivalently $\frac{2\ell}{n} < \ell^2 - 2$. Just choose $n > \frac{2\ell}{\ell^2 - 2}$.

Remark

If $S = \emptyset$, then vacuously any real number is an upper bound.

Proposition 2.9 *Let S be a non-empty subset of \mathbf{R} which is bounded above. The $\ell = \sup S$ iff:*

1. ℓ is an upper bound for S .
2. given any (real) $\epsilon > 0$, $\exists s \in S$ with $s > \ell - \epsilon$.

Proof Suppose first $\ell = \sup S$. Then (1) is true by definition. Now for any $\epsilon > 0$, since ℓ is the least upper bound, $\ell - \epsilon$ is not an upper bound for S . So $\exists s \in S$ such that $s > \ell - \epsilon$.

Conversely, suppose (1) and (2) hold. Then by (1) ℓ is an upper bound. We want to see that it's the least such. So let U be any upper bound for S . We want $U \geq \ell$. Suppose for a contradiction $U < \ell$. Put $\epsilon = \ell - U$ in (2). So $\exists s \in S$ with $s > \ell - \epsilon = \ell - (\ell - U) = U$. So U is not an upper bound.

Proposition 2.10 *Suppose $S, T \subseteq \mathbf{R}$, non-empty and bounded above. Then $S \cup T$ is bounded above and $\sup(S \cup T) = \max(\sup S, \sup T)$.*

Proof: exercise.

Proposition 2.11 *Suppose S, T as in 2.10. Let $S + T = \{x + y \in \mathbf{R} : x \in S, y \in T\}$. Then $S + T$ is bounded and $\sup(S + T) = \sup S + \sup T$.*

Proof For any $x \in S, y \in T$ we have $x \leq \sup S, y \leq \sup T \Rightarrow \sup S + \sup T$ is an upper bound for $S + T$.

Let $\epsilon > 0$. BY 2.9, $\exists s \in S$ such that $s > \sup S - \frac{\epsilon}{2}$ and $\exists t \in T$ such that $t > \sup T - \frac{\epsilon}{2}$. Then $s + t \in S + T$, and $s + t > \sup S + \sup T - \epsilon$. So by 2.9, $\sup S + \sup T = \sup(S + T)$.

Warning: Given $\ell = \sup S$ and $\epsilon > 0$, it's not necessarily true that $\ell - \epsilon \in S$. All you can be sure of is that $\exists s \in S$ with $s > \ell - \epsilon$.

Example: $S = \{1 - \frac{1}{n} : n \in \mathbf{N}_+\}$. Then $\sup S = 1$, and $\epsilon = 3/10 > 0$, $\ell - \epsilon = 7/10$, not in S , but for example $9/10 \in S$.

Definition 2.12 *Given $S \subseteq \mathbf{R}$, non-empty and bounded below, a greatest lower bound G for S is a number such that:*

1. g is a lower bound of S .
2. g is the greatest such.

Written $g = \inf S$ (infimum).

Proposition 2.13 *If S is non-empty, $S \subseteq \mathbf{R}$, and bounded below, then $\exists \inf S$.*

Proof follows from Completeness Axiom by a trick: let $S^* = \{x \in \mathbf{R} : -x \in S\}$. Easy to check: x is a lower bound for S iff $-x$ is an upper bound for S^* .

Outline: Since S is bounded below, S^* is bounded above and non-empty, so by Completeness $\sup S^*$ exists. Not hard to check that $-\sup S^*$ is a greatest lower bound for S .

3 Revision on limits of functions

Definition 3.1 Suppose a real-valued function f is defined at least 'near' a , i.e. $\exists(b, c)$ with $a \in (b, c)$ and f defined at least on $(b, c) \setminus \{a\}$. Then $\lim_{x \rightarrow a} f(x) = \ell$ means: Given any $\epsilon > 0$, $\exists \delta > 0$ such that whenever $0 < |x - a| < \delta$, we have $|f(x) - \ell| < \epsilon$.

N.B. For $\lim_{x \rightarrow a} f(x)$ to exist, it's not necessary for f to be defined at a , and if $f(a)$ is defined it need not be equal to $\lim_{x \rightarrow a} f(x)$.

Example

Suppose $f : [0, 1] \rightarrow \mathbf{R}$ is defined by

$$f(x) = \begin{cases} x & : x \neq 1/2 \\ 666 & : x = 1/2 \end{cases}$$

Then $\lim_{x \rightarrow 1/2} f(x) = 1/2 \neq f(1/2)$. (Or: don't define $f(1/2)$.)

Example

$$\text{Let } f(x) = \begin{cases} \sin(1/x) & : x \neq 0 \\ 0 & : x = 0 \end{cases}$$

Want to show that $\lim_{x \rightarrow 0} f(x)$ does not exist.

Three levels:

1. (intuitive) The graph oscillates too much near 0 for $\lim_{x \rightarrow 0} f(x)$ to exist.
2. (towards proof) Take any potential limit ℓ . Then near 0, there are points where the graph is at +1 and points where it's at -1. These can't both be arbitrarily close to ℓ .
3. Let $\ell \in \mathbf{R}$ and take $\epsilon = 1/2$, say. Want to show that 3.1 doesn't work for this f, ϵ, ℓ .

Take any $\delta > 0$. Let $N \in \mathbf{Z}$, $N > 1/\delta$ and consider

$$x_1 = \frac{1}{(2N + \frac{1}{2})\pi}, \quad x_2 = \frac{1}{(2N + \frac{3}{2})\pi}.$$

We see $f(x_1) = 1$, $f(x_2) = -1$, and $0 < |x_1 - 0| < \delta$. Now if $0 < |f(x) - \ell| < 1/2$ whenever $0 < |x - 0| < \delta$, we'd have $|1 - \ell| < 1/2$, $|-1 - \ell| < 1/2$ so $2 = |1 - (-1)| = |1 - \ell + -1 - \ell| \leq |1 - \ell| + |-1 - \ell| < 1$. Contradiction.

One-sided limits

If f is defined at least on (a, b) ($b > a$) we say $\lim_{x \rightarrow a^+} f(x) = \ell$ or $f(x) \rightarrow \ell$ as $x \rightarrow a^+$ if given any $\epsilon > 0$, $\exists \delta > 0$ such that $|f(x) - \ell| < \epsilon$ whenever $a < x < a + \delta$.

Similarly for $\lim_{x \rightarrow b^-} f(x)$, $\lim_{x \rightarrow \infty} f(x)$, $\lim_{x \rightarrow -\infty} f(x)$.

Definition 3.2 Suppose $f : (b, c) \rightarrow \mathbf{R}$ and $a \in (b, c)$. We say f is continuous at a if $f(x) \rightarrow f(a)$ as $x \rightarrow a$.

Variants: $f : [a, b] \rightarrow \mathbf{R}$ is (right-hand) continuous at a is $f(x) \rightarrow f(a)$ as $x \rightarrow a+$. Similarly for left-hand continuity at b .

Proposition 3.3 $f : (b, c) \rightarrow \mathbf{R}$ is continuous at $a \in (b, c)$ iff given any $\epsilon > 0$, $\exists \delta > 0$ such that $|f(x) - f(a)| < \epsilon$ whenever $|x - a| < \delta$.

Proof follows from definitions 3.1 and 3.2.

Note Difference from 3.1: no “ $0 < |x - a|$ ” this time. When $x = a$, $f(x) = f(a)$ so $|f(x) - f(a)| = 0 < \epsilon$, i.e. this time we’re looking at $f(x) \rightarrow f(a)$ as $x \rightarrow a$, not $f(x) \rightarrow \ell$ as $x \rightarrow a$.

There are two ways of thinking about continuity of f at a :

1. Approximation: think of f as a black box which accepts real number inputs and gives you real number outputs. Then if you vary input x only slightly from a , output $f(x)$ also varies only slightly from $f(a)$.
2. Geometrically: graph doesn’t jump around too much: given any horizontal strip centred on height $f(a)$ you can find a vertical strip centred on a , narrow enough so that all of the graph in the vertical strip is also inside the horizontal strip. For some reason he called this the “principle of inertia”.

Definition 3.4 $f : [a, b] \rightarrow \mathbf{R}$ is said to be continuous if it’s continuous at every point in $[a, b]$ (RH continuous at a , LH continuous at b).

Note if $f : [a, b] \rightarrow \mathbf{R}$ is continuous and I is any interval $I \subseteq [a, b]$, then $f|_I : I \rightarrow \mathbf{R}$ is also continuous (obvious if you ask me). For $c \in I$, any $\epsilon > 0$, by continuity of f at c , $\exists \delta > 0$ such that $|f(x) - f(c)| < \epsilon$ whenever $|x - c| < \delta$ and $x \in [a, b]$ so certainly true $\forall x$ with $|x - c| < \delta$ and $x \in I$.

Continuous functions behave well under addition and multiplication.

Proposition 3.5 Suppose $f, g : (b, c) \rightarrow \mathbf{R}$ and $a \in (b, c)$ and f, g continuous at a . Then so are the following functions:

1. $x \mapsto kf(x)$ (k constant)
2. $x \mapsto f(x) + g(x)$ (“ $f + g$ ”)
3. $x \mapsto f(x)g(x)$
4. provided $g(a) \neq 0$, $x \mapsto 1/g(x)$

Proof: Given definition 3.2, follows immediately from Algebra of Limits.

Corollary 3.6 Any polynomial $p(x)$, $x \mapsto a_0 + a_1x + \dots + a_nx^n$, a_i constant, is continuous at any point in \mathbf{R} and so is any rational function $x \mapsto \frac{p(x)}{q(x)}$ where p and q are polynomials, except at (finite set of) points where $q(x) = 0$.

Proof A sequence of easy inductions. But first,

1. Any constant function $x \mapsto k$ is continuous, since for any $\epsilon > 0$ take any $\delta > 0$ and we have $|k - k| = 0 < \epsilon$ whenever $|x - a| < \delta$.
2. The function $x \mapsto x$ is continuous, since for any $\epsilon > 0$, take $\delta = \epsilon$. Then $|x - a| < \delta \Rightarrow |f(x) - f(a)| = |x - a| < \epsilon$.

First induction: $\forall n \in \mathbf{N}$, $x \mapsto x^n$ is continuous. Anchor is (2) above. Induction step follows from $x \mapsto x$ and $x \mapsto x^n$ continuous $\Rightarrow x \mapsto x^{n+1}$ continuous.

Then by one more application of ‘product’, $x \mapsto a_n x^n$ is continuous. Finally $x \mapsto a_0 + a_1 x + \dots + a_n x^n$ is continuous by easy induction on adding two continuous functions.

Also, $p(x)/q(x)$ follows by “quotient” part of Algebra of Limits.

Remark Continuity of f at a can be split into four bits:

1. $\lim_{x \rightarrow a+}$ exists.
2. $\lim_{x \rightarrow a+} = f(a)$
3. $\lim_{x \rightarrow a-}$ exists.
4. $\lim_{x \rightarrow a-} = f(a)$

Examples

$$f(x) = \begin{cases} 0 & : x \leq 0 \\ 1 & : x > 0 \end{cases}, a = 0. \text{ 1, 3, 4 true but 2 fails at 0.}$$

$$f(x) = \begin{cases} 0 & : x < 0 \\ 1 & : x \geq 0 \end{cases}, a = 0. \text{ 1, 2, 3 hold but 4 fails at 0.}$$

These two functions have a *simple jump discontinuity*.

$$f(x) = \begin{cases} \sin(1/x) & : x > 0 \\ 0 & : x \leq 0 \end{cases}, a = 0. \text{ 3, 4 true, 1 fails so 2 meaningless.}$$

Similarly we can get 1 and 2 to hold and 3 to fail so 4 is meaningless.

Proposition 3.7 *The composition $f \circ g$ of continuous functions is continuous.*

Proposition 3.8 *If f is differentiable at $a \in \mathbf{R}$, it's continuous at a .*

$$\frac{f(x) - f(a)}{x - a} \rightarrow f'(a) \text{ as } x \rightarrow a \Rightarrow \underbrace{f(x) - f(a)}_{\text{hence also } \rightarrow 0} - \underbrace{(x - a)f'(a)}_{\text{must tend to 0}} \rightarrow 0 \text{ as } x \rightarrow a$$

4 Results about sequences proved using Completeness, and needed to study Continuity

Proposition 4.1 *Any monotone bounded sequence converges.*

Proof First consider (a_n) increasing, bounded above. Let $S = \{a_n : n \in \mathbf{N}\}$. Certainly non-empty, and bounded above. So the least upper bound $\sup S$ exists — call it ℓ . We'll show (a_n) converges to ℓ .

Let $\epsilon > 0$. The $\ell - \epsilon$ is not an upper bound for S . So $\exists N$ such that $a_N < \ell - \epsilon$. By an easy induction using the fact $a_{n+1} \geq a_n \forall n$, get $a_n > \ell - \epsilon \forall n \geq N$. But also $a_n \leq \ell \forall n$. So $\forall n \geq N$ we have $\ell - \epsilon < a_n \leq \ell$, so $|a_n - \ell| < \epsilon$.

Finally, if (a_n) is decreasing and bounded below, then $(-a_n)$ is increasing and bounded above, so $-a_n \rightarrow \ell$, say, as $n \rightarrow \infty$, so it's easily seen that $a_n \rightarrow -\ell$ as $n \rightarrow \infty$.

Theorem 4.2 (Bolzano-Weierstrass) *Any bounded sequence of real numbers has at least one convergent subsequence.*

Proof I It is enough to show that (a_n) has a monotonic subsequence. For then that subsequence is monotonic and bounded so 4.1 applies. We can show ...

Theorem 4.3 *Any sequence of real numbers has a monotonic subsequence.*

Local definition: for a sequence (a_n) , n is a scenic viewpoint if $\forall m \geq n, a_n \geq a_m$.

Proof The sequence (a_n) has infinitely many scenic viewpoints, or it doesn't.

Case 1 Suppose that there are infinitely many scenic viewpoints. Put them in order: $n_1 < n_2 < n_3 < \dots$. Consider a_{n_r} : since n_r is a scenic viewpoint, $a_m \leq a_{n_r} \forall m \geq n_r$, in particular $a_{n_{r+1}} \leq a_{n_r}$. This is true $\forall r$ so (a_{n_r}) is decreasing.

Case 2 There are not infinitely many scenic viewpoints. Let n_0 be the largest scenic viewpoint, if any exist, otherwise let $n_0 = 0$. Let $n_1 = n_0 + 1$. Then n_1 is not a scenic viewpoints so $\exists m > n_1$ such that $a_m > a_{n_1}$. Choose any such m and call it n_2 . Inductively: suppose $n_1 < n_2 < \dots < n_r$ chosen so that $a_{n_1} < a_{n_2} < \dots < a_{n_r}$. Then n_r is not a scenic viewpoint, so we can choose $n_{r+1} > n_r$ with $a_{n_{r+1}} > a_{n_r}$. This way, inductively, we get a (strictly) increasing subsequence of (a_n) .

Proof II (Bisection method)

Definition: call an interval I crowded if $a_n \in I$ for infinitely many n .

Step 1. (a_n) is bounded so $\exists K \in \mathbf{R}$ such that $a_n \in [-K, K] \forall n$. So either $\dagger[-K, 0]$ is crowded or $[0, K]$ is crowded. If \dagger holds, put $x_0 = -K, y_0 = 0$. If \dagger fails, put $x_0 = 0, y_0 = K$. Then $[x_0, y_0]$ is crowded.

Suppose inductively we've already chosen $x_0 \leq x_1 \leq \dots \leq x_r \leq y_r \leq y_{r-1} \leq \dots \leq y_0$ such that $[x_r, y_r]$ is crowded and $y_i - x_i = K/2^i$ for $i = 0, 1, \dots, r$.

Bisect $[x_r, y_r]$. Either \ddagger the left-hand half or the right-hand half is crowded. If \ddagger holds, put $x_{r+1} = x_r, y_{r+1} = (x_r + y_r)/2$. Otherwise put $x_{r+1} = (x_r + y_r)/2, y_{r+1} = y_r$. Then clearly the assumption hold for $n + 1$.

Not (x_r) in increasing and bounded above (by any y_s) so $x_r \rightarrow x$ as $r \rightarrow \infty$ and similarly $y_r \rightarrow y$ as $r \rightarrow \infty$ and at any stage r we have $x_r \leq x \leq y \leq y_r$ ($y_s \geq x$ for any s since $x_r \leq y_s$ (fixed) $\forall r$). Put $y_r - x_r = K/2^r$ — true $\forall r$ so by sandwiching, $y = x$.

Now need to see that a subsequence of (a_n) converges to x . Every $[x_r, y_r]$ is crowded so $a_{n_r} \in [x_r, y_r]$ for some n_r . But this is not yet an orderly subsequence ($n_{r+1} > n_r \forall r$). Suppose we already have $n_1 < n_2 < \dots < n_r$ such that $a_{n_i} \in [x_i, y_i] \forall i \leq r$. Since $[x_{r+1}, y_{r+1}]$ is crowded, we can choose $n_{r+1} > n_r$ such that $a_{n_{r+1}} \in [x_{r+1}, y_{r+1}]$. Now we see $x_r \leq a_{n_r} \leq y_r \forall r$, so $a_{n_r} \rightarrow x$ as $r \rightarrow \infty$, by sandwiching.

There's another version of B-W — see e.g. Spivak *Calculus* p. 386 ex. 21.

5 Results about continuous functions $f : [a, b] \rightarrow \mathbf{R}$

Theorem 5.1 *Any continuous function $f : [a, b] \rightarrow \mathbf{R}$ is bounded and attains its bounds.*

We now rephrase the second part of this assertion thrice.

Suppose $f : [a, b] \rightarrow \mathbf{R}$ is bounded, so $\exists K$ such that $|f(x)| \leq K \forall x \in [a, b]$. Equivalently, $f([a, b]) = \{y \in \mathbf{R} \text{ such that, for some } x \in [a, b], f(x) = y\} \subseteq [-K, K]$.

So image $f([a, b])$ is (non-empty) bounded above and below, so has a sup and an inf; (ii) says that they are taken on as values of f . Equivabloodily, $\exists c, d \in [a, b]$ s.t. $f(c) \leq f(x) \leq f(d) \forall x \in [a, b]$.

Remark It's crucial that $[a, b]$ is closed and bounded.

Examples

- Define $f : (0, 1) \rightarrow \mathbf{R}$ by $f(x) = 1/x$. f continuous, not bounded.
- Define $g : (0, 1) \rightarrow \mathbf{R}$ by $g(x) = x$. Continuous, bounded but bounds not attained.
- Define $g : [1, \infty) \rightarrow \mathbf{R}$ by $g(x) = x$. Not bounded.
- Define $f : [1, \infty) \rightarrow \mathbf{R}$ by $f(x) = 1/x$. Bounded but infimum not attained.

Recall If (a_n) converges to ℓ and $a_n \geq L \forall n$, then $\ell \geq L$ (also, $a_n > L \Rightarrow \ell \geq L$, not $\ell > L$). Similarly, $a_n \leq L \forall n \Rightarrow \ell \leq L$.

5.0.1 Proof

(i) Suppose f not bounded. Then for any $n \in \mathbf{N}$, $\exists a_n \in [a, b]$ with $|f(a_n)| > n$. Consider (a_n) — bounded, so by Bolzano-Weierstrass, has a convergent subsequence a_{n_r} converging to some $\ell \in \mathbf{R}$ †.

Now $a \leq a_{n_r} \leq b \forall r$, so $a \leq \ell \leq b$, i.e. $\ell \in [a, b]$ ‡. Use continuity of f at ℓ : get $f(a_{n_r}) \rightarrow f(\ell)$ as $r \rightarrow \infty$. All we need of this is: $(f(a_{n_r}))$ converges, hence (from result last term) is bounded. But $|f(a_{n_r})| > n_r \forall r$. Contradiction, so f bounded.

†is where we use $[a, b]$ bounded. ‡is where we use $[a, b]$ closed.

(ii) Let g, ℓ be the inf and sup of $f([a, b])$. Suppose e.g. ℓ is not attained on $[a, b]$. We know $f(x) \leq \ell \forall x \in [a, b]$ and by our supposition $f(x) < \ell \forall x \in [a, b]$. **KEY:** Let $h(x) = 1/(\ell - f(x))$, $x \in [a, b]$. By (3.5), h is continuous on $[a, b]$. By (1), h is bounded on $[a, b]$. Say

$1/(\ell - f(x)) \leq U \Rightarrow \ell - f(x) \geq 1/U \Rightarrow f(x) \leq \ell - 1/U \forall x \in [a, b]$. i.e. $\ell - 1/U$ is an upper bound for $f([a, b])$ — contradicts leastness of ℓ . Proof that $\inf f([a, b])$ is attained on $[a, b]$ is similar.

Theorem 5.2 (Special case of IVT) Suppose $f : [a, b] \rightarrow \mathbf{R}$ is continuous, and $f(a) < 0$ and $f(b) > 0$. Then $\exists c \in (a, b)$ with $f(c) = 0$.

Note not true for **Q**: try $f : [0, 2] \rightarrow \mathbf{R}$, $f(x) = x^2 - 2$.

Proof I

Consider $s = \{x \in [a, b] : f(x) \leq 0\}$. Note $s \neq \emptyset$ ($a \in S$) and bounded above (e.g. by b). So by Completeness, $\exists \sup S$. Let $c = \sup S$. We'll show $f(c) = 0$ by contradicting the other possibilities.

1. Suppose $f(c) > 0$. We'll show this contradicts c being *least* upper bound of S . By continuity of f at c , $\exists \delta > 0$ such that whenever $x \in [a, b]$ and $|x - c| < \delta$, we have $|f(x) - f(c)| < f(c)/2$ (note $f(c)/2 > 0$).

For all such x , $-(f(x) - f(c)) < f(c)/2$, so $f(x) > f(c) - f(c)/2 > 0$. Now $c \neq 0$ (we're supposing $f(c) > 0$, and we know $f(a) < 0$). We may as well assume $\delta < c - a$, so $c - \delta > a$. So $\forall x \in (c - \delta, c]$ we have $f(x) > 0$. By definition of S , $f(x) > 0 \forall x > c$ (in $[a, b]$). So $c - \delta$ is an upper bound for S , contradicting leastness of c .

2. Suppose $f(c) < 0$. We'll contradict " c is an upper bound for S ". By continuity of f at c , $\exists \delta > 0$ such that whenever $x \in [a, b]$ and $|x - c| < \delta$, we have $|f(x) - f(c)| < -f(c)/2$ (note $-f(c)/2 > 0$).

So for all such x we have $f(x) - f(c) < -f(c)/2$ so $f(x) < f(c) - f(c)/2 < 0$. Now $c \neq b$ ($f(c) < 0$, $f(b) > 0$). So certainly $\exists x \in [a, b]$ with $c < x < c + \delta$, and for any one such, $f(x) < 0$, so contradicting " c an upper bound for S ".

So $f(c) = 0$.

Proof II

Proceed similarly to Bolzano-Weierstrass proof, but choose $a = x_0 \leq x_1 \leq \dots \leq x_r \leq y_r \leq y_{r-1} \leq \dots \leq y_0 \leq b$ as follows:

If $f((a+b)/2) \leq 0$, put $x_1 = (a+b)/2$, $y_1 = b$.

If $f((a+b)/2) > 0$, put $x_1 = a$, $y_1 = (a+b)/2$.

Continue inductively — get $a \leq x_1 \leq \dots \leq x_r \leq y_r \leq y_{r-1} \leq \dots \leq y_1 \leq b$. such that $y_r - x_r = (b-a)/2^r$ and $f(x_r) \leq 0$, $f(y_r) \geq 0$. As in Bolzano-Weierstrass, $(x_r), (y_r)$ converge to a common limit c and $f(c) \leq 0$ by continuity of f at c , and $f(c) \geq 0$ ($f(x_r) \leq 0 \forall r$) so $f(c) = 0$.